

Recall: The simplex method works under the following assumptions:

- 1) The program is in the equality form: we want to maximize

$$Z = c_1x_1 + \dots + c_nx_n$$

subject to the constraints

$$\begin{aligned} a_{11}x_1 + \dots + a_{1n}x_n &= b_1 \\ \dots \quad \dots \quad \dots \quad \dots \quad \dots \\ a_{m1}x_1 + \dots + a_{mn}x_n &= b_m \\ x_1, x_2, \dots, x_n &\geq 0 \end{aligned}$$

- 2) The coefficient matrix

$$A = \begin{bmatrix} a_{11} & \dots & a_{1n} \\ \vdots & & \vdots \\ a_{m1} & \dots & a_{mn} \end{bmatrix}$$

is in the basic form.

- 3) $b_i \geq 0$ for $i = 1, \dots, m$.

- These assumptions assure that the program has a basic feasible solution.
- **Phase I of the simplex method** finds some basic feasible solution or verifies that no feasible solutions exists.

Example. Maximize

$$z = x_1 + 2x_2$$

subject to:

$$x_1 + 3x_2 + x_3 = 2$$

$$-x_2 - x_3 = -1$$

$$x_1, x_2, x_3 \geq 0$$

To get started we need some basic feasible solution

Attempt 1: Use x_1, x_2 as basic variables:

$$(-1) \cdot \begin{bmatrix} x_1 & x_2 & x_3 \\ \boxed{1} & 3 & 1 & | & 2 \\ 0 & \boxed{-1} & -1 & | & -1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 3 & 1 & | & 2 \\ 0 & 1 & 1 & | & 1 \end{bmatrix} \xrightarrow{(-3)} \begin{bmatrix} \boxed{1} & 0 & -2 & | & -1 \\ 0 & \boxed{1} & 1 & | & 1 \end{bmatrix}$$

Solution:

$$\begin{array}{l|l} \text{free} & \begin{cases} x_3 = 0 \\ x_1 = -1 \\ x_2 = 1 \end{cases} \\ \text{basic} & \end{array}$$

Not feasible since $x_1 < 0$.

Attempt 2: Use x_1 and x_3 as free variables.

$$(-1) \cdot \begin{bmatrix} x_1 & x_2 & x_3 \\ \boxed{1} & 3 & 1 & | & 2 \\ 0 & -1 & \boxed{-1} & | & -1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 3 & 1 & | & 2 \\ 0 & 1 & 1 & | & 1 \end{bmatrix} \xrightarrow{(-1)} \begin{bmatrix} \boxed{1} & 3 & 0 & | & 1 \\ 0 & 1 & \boxed{1} & | & 1 \end{bmatrix}$$

Solution:

$$\begin{array}{l|l} \text{free} & \begin{cases} x_3 = 0 \\ x_1 = 1 \\ x_2 = 1 \end{cases} \\ \text{basic} & \end{array}$$

This is a basic feasible solution, so we can use it to get the simplex method started.

Phase I of the simplex method

Assumption: The linear program is in the equality form with constraints

$$a_{i1}x_1 + \dots + a_{in}x_n = b_i$$

for $i = 1, \dots, m$.

1) Modify the constraints, if needed, so that $b_i \geq 0$ for $i = 1, \dots, m$.

$$\begin{array}{l} x_1 + 3x_2 + x_3 = 2 \\ (-1) \cdot (-x_2 - x_3 = -1) \end{array} \quad \rightarrow \quad \begin{array}{l} x_1 + 3x_2 + x_3 = 2 \\ x_2 + x_3 = 1 \end{array}$$

2) Add an additional variable s_i to the i -th constraint for $i = 1, \dots, m$.

$$\begin{array}{l} x_1 + 3x_2 + x_3 = 2 \\ x_2 + x_3 = 1 \end{array} \quad \rightarrow \quad \begin{array}{l} x_1 + 3x_2 + x_3 + s_1 = 2 \\ x_2 + x_3 + s_2 = 1 \end{array}$$

Note. The augmented matrix of the new constraints is in the basic form.

$$\left[\begin{array}{ccc|cc|c} x_1 & x_2 & x_3 & s_1 & s_2 & \\ \hline 1 & 3 & 1 & 1 & 0 & 2 \\ 0 & 1 & 1 & 0 & 1 & 1 \end{array} \right]$$

free basic

3) Use the simplex method to minimize the function

$$Z = s_1 + \dots + s_m$$

with the new constraints. If the minimum is non-zero, then the original linear program has not feasible solutions.

If the minimum is $z = 0$, then the solution that gives the minimum has $s_i = 0$ for $i = 1, \dots, m$. In such case, values of the variables x_1, \dots, x_n give a basic feasible solution of the original linear program.